## Arc segmentation in linear time

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## Contribution

We propose a linear algorithm based on discrete geometry approach for segmentation of a curve into digital arcs

## Motivation

- Arc and circle are basic object in discrete geometry
- Arc and circle appear often in images
- Shape contains often digital arcs.
$\Rightarrow$ The study of these primitives is important



## Discrete circle

- Basic object in discrete geometry
- Based on discritization of a real circle


## [Kim84] <br> [Nakamura84] <br> [Andres04]

## Discrete line [Reveillès91]

A discrete line, noted as $D(a, b, \mu, \omega), a, b, \mu, \omega \in \mathbb{Z}^{2}$ is a set of points that verifies: $\mu \leq a x-b y<\mu+\omega$

## Arc in the tangent space

Proposition 1:
Let $C=\left\{C_{i}\right\}_{i=0}^{n}$ be a polygon, $\alpha_{i}=$ $\angle\left(C_{i-1} C_{i}, C_{i} C_{i+1}\right)$. The length of $C_{i} C_{i+1}$ is $l_{i}$, for $i \in\{0, \ldots, n-1\}$. The vertices of $C$ are on a real arc of radius $R$ with center $O, \angle C_{i} O C_{i+1} \leq \frac{\pi}{4}$ for $i \in\{1, \ldots, n-1\}$. This results below is obtained.
$\frac{1}{R}<\frac{\alpha_{i}}{\frac{l_{i}+l_{i+1}}{2}}<\frac{1}{0.9742979 R}$

## Interest of proposition 1



Example:


## Deciding if a curve is an arc

1. Polygonalize the input digital curve by polygon $P$ based on recognition of BS of width 1
2. Transform $P$ to $T(P)$ in the tangent space
3. Determine the midpoint set $M p C=\left\{M_{i}\right\}_{i=1}^{n-1}$ of horizontal segment of $T(P)$
4. Verify if $M p C$ is a BS of width $\epsilon$ [Debled 06]

## Arc segmentation

## Main ideas:

1. Polygonalize the input curve
2. Transform the polygon to tangent space
3. Construct the curve of midpoints in the tangent space
4. Polygonalize the midpoint curve - Utilize parameter $\alpha$ to verify detected arcs

## Study of quasi co-linear property

Convergence of radius of local circumcircles Proposition 2:
Let $C \xrightarrow{=}\left\{C_{i}\right\}_{i=0}^{n}$ be a polygon, $\alpha_{i}=$ $\angle\left(\overrightarrow{C_{i-1} C_{i}}, \overrightarrow{C_{i} C_{i+1}}\right)$. The length of $C_{i} C_{i+1}$ is $l_{i}$, for $i \in\{0, \ldots, n-1\}$. We denote $O_{i}, R_{i}, H_{i}$ respectively the center and the radius of circumcirle that passes to 3 points $C_{i-1}, C_{i}, C_{i+1}$, the projection of $O_{i}$ on $C_{i} C_{i+1}$, suppose that $R_{i}-O H_{i} \leq h$ for $i \in\{1, \ldots, n-1\}$. This results below is obtained. $R_{i} \alpha_{i} \geq \frac{l_{i-1}+l_{i}}{2} \geq R_{i} \alpha_{i}-0.3377 h \alpha_{i}$
Convergence of centers of local circumcircles Proposition 2:
Let us consider a sequence of points $\{C\}_{i=0}^{n}$. We denote $O_{i}^{\prime}$ (resp. $O_{i}$ ) and $R_{i}$ (resp. $R_{i}$ ) are the center and radius of circumcircle that passes to 3 points $C_{0}$ (resp. $C_{1}$ ), $C_{i}, C_{i+1}$. There exist $R$ and $\delta$ satisfied $R, \delta \in \mathcal{R}, 0 \leq R_{i}-R \leq \delta, i=1, \ldots, n-1$. Suppose that $\angle C_{k} C_{j} C_{j+1}>\frac{\pi}{2}$ for $k \in\{0,1\}, k<$ $j<n$. Therefore, we have this property $0 \leq R_{i}^{\prime}-$ $R \leq \delta, 0 \leq R_{i}^{\prime \prime}-R \leq \delta$, for $1 \leq i \leq n-1$.


## Experimentation



